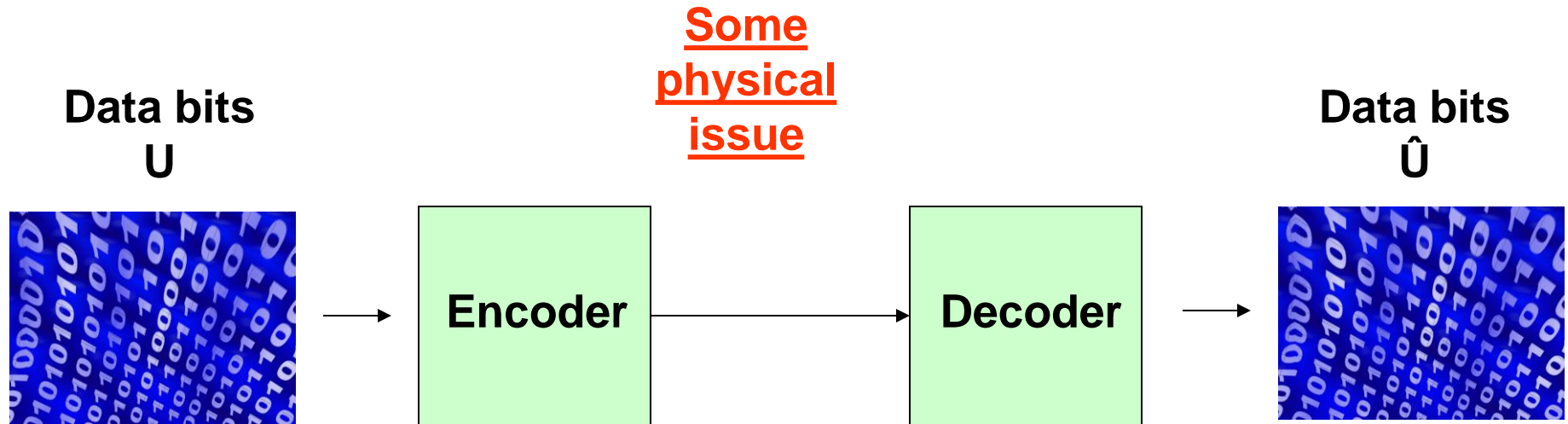




Seminar on Coding for Non-Volatile Memories  
236803/048704 – CS/EE Departments, Technion

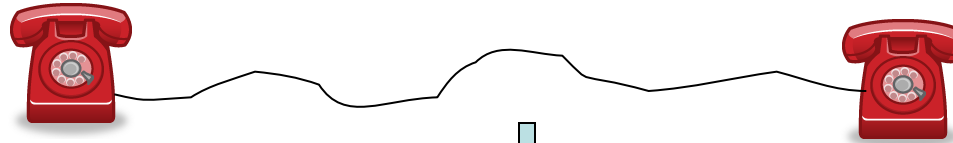
# **ASYMMETRIC, LIMITED-MAGNITUDE ERRORS**

# Coding: The General Problem



Find good codes that overcome physical issue and give  $U = \hat{U}$

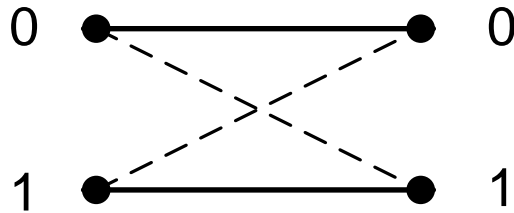
# Classical Coding Theory



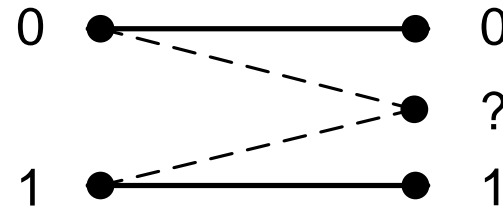
Shannon



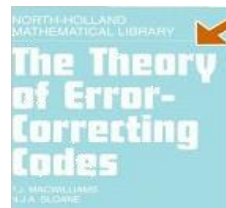
BSC



BEC

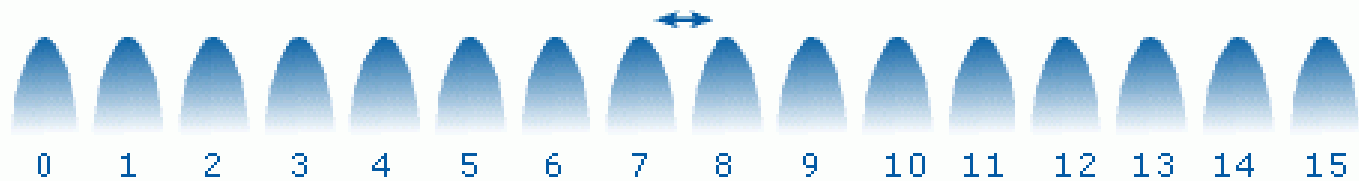
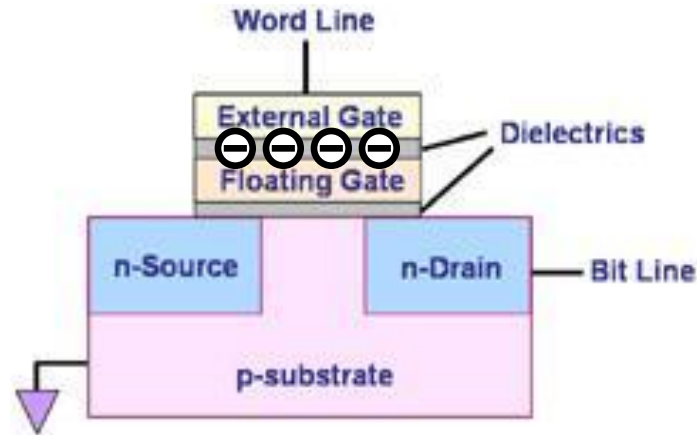


Hamming  
(distance)

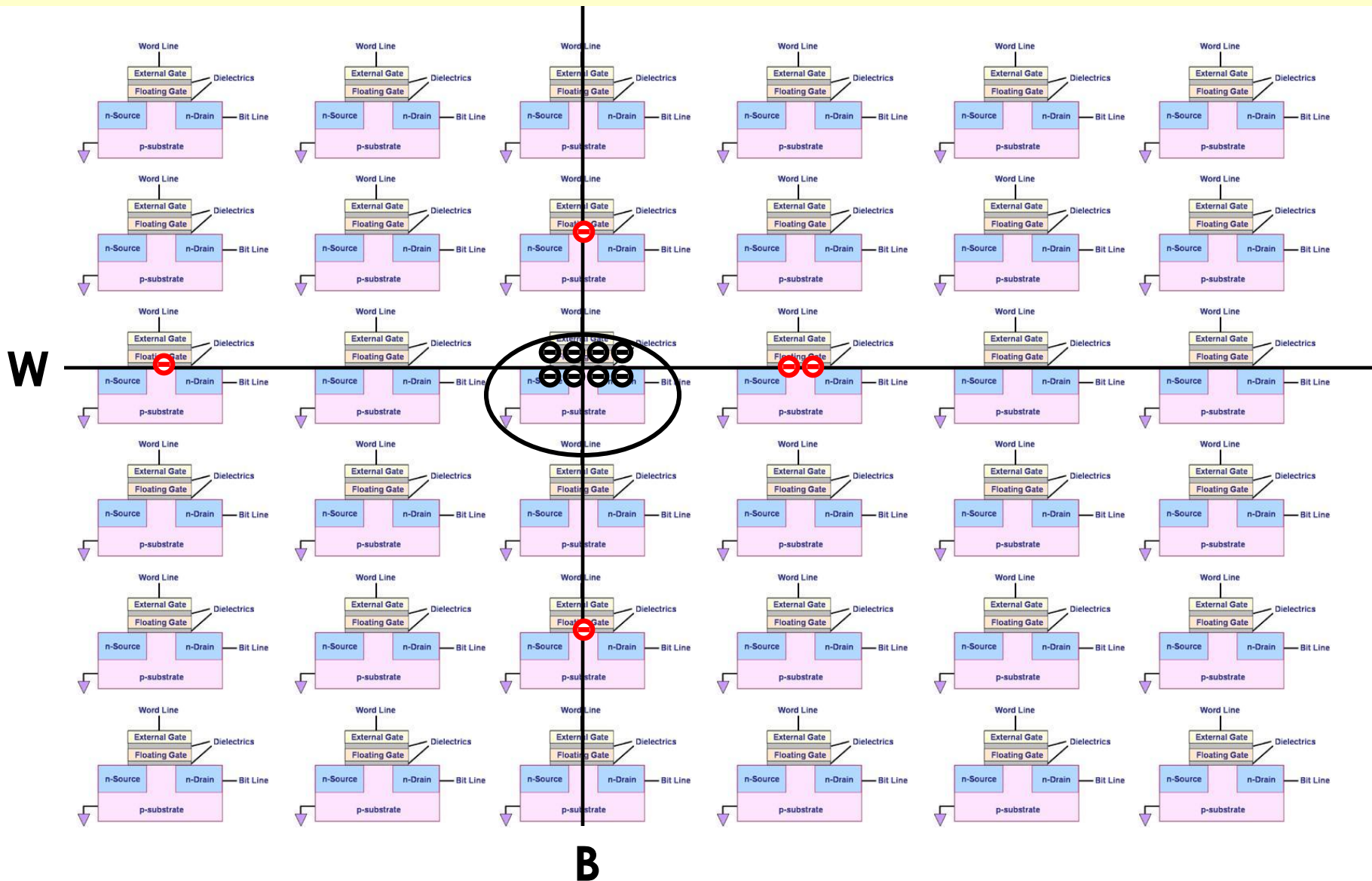


Codes for **Symmetric Symbol** Errors

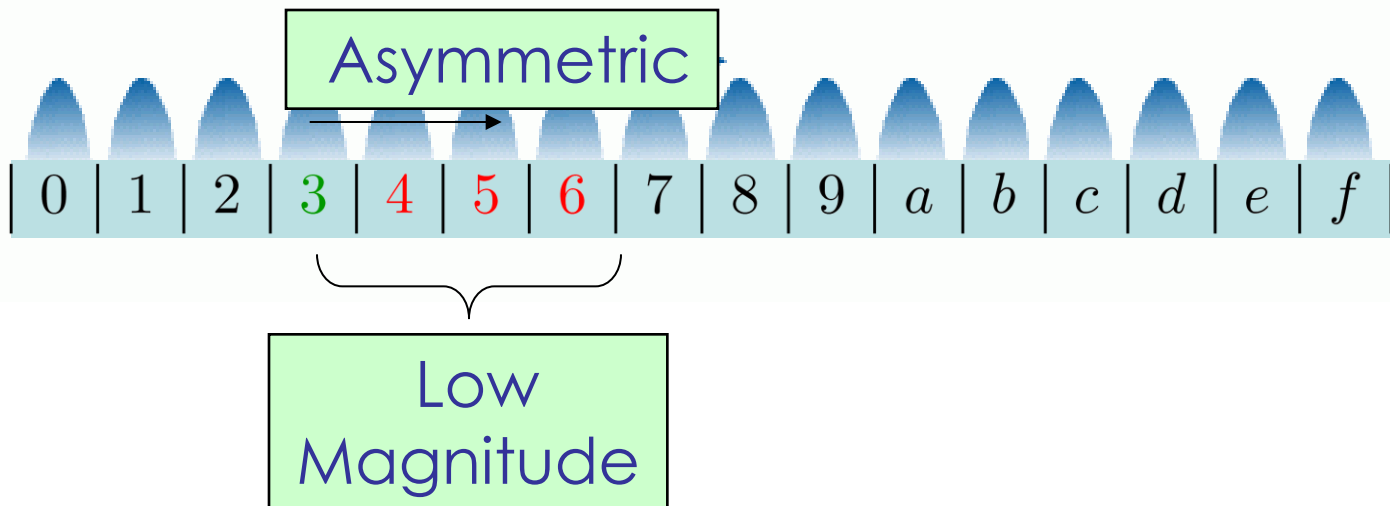
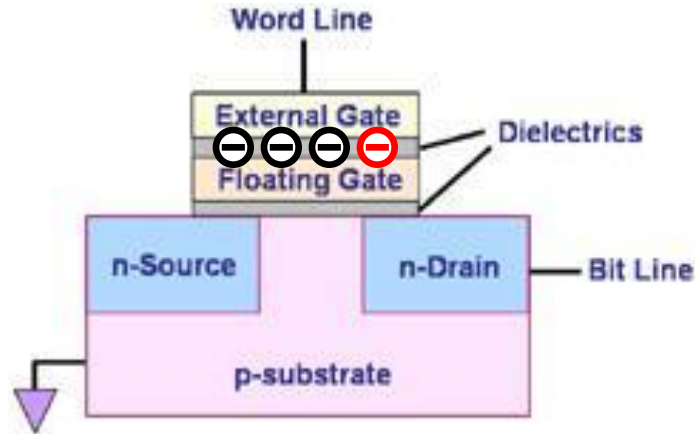
# Multi-Level Flash Storage



# Dominant Flash Errors



# Multi-Level Flash Errors



# Asymmetric Limited-Magnitude Errors

Stored:  $\vec{c} = [c_1, c_2, \dots, c_n] \in Q^n$

Corrupted:  $\vec{r} = [r_1, r_2, \dots, r_n] \in Q^n$

**A**symmetric  
**l**imited-  
**M**agnitude  
Errors

$$0 \leq r_i - c_i \leq \ell, \quad w_H(\vec{r} - \vec{c}) \leq t$$

Research Problem:

Code  $C \subset Q^n$  that corrects  $t$  *ALM* errors

For each  $(n, \ell, t)$

# Example, $t = 1, \ell = 1$

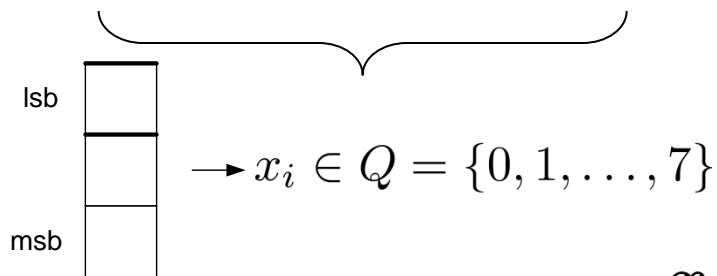
18 information bits

$n = 7$  8-ary cells ( $q=8$ )

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 0 | 1 | 1 | 0 |   |   |   |
| 0 | 0 | 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 0 | 1 | 1 |

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 0 | 1 | 1 | 0 | 0 | 1 | 1 |
| 0 | 0 | 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 0 | 1 | 1 |

$\in \Sigma_H$ , the Binary Hamming Code



$$\mathbf{x} = \boxed{4} \boxed{5} \boxed{3} \boxed{6} \boxed{2} \boxed{5} \boxed{7} \in C$$



# Decoding Example

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 4 | 5 | 3 | 6 | 2 | 5 | 7 |
|---|---|---|---|---|---|---|

A1M ↓

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 4 | 5 | 4 | 6 | 2 | 5 | 7 |
|---|---|---|---|---|---|---|



|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 0 | 1 | 0 | 0 | 0 | 1 | 1 |
| 0 | 0 | 0 | 1 | 1 | 0 | 1 |
| 1 | 1 | 1 | 1 | 0 | 1 | 1 |

\*

Decode  $\Sigma$

-1

+1

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 4 | 5 | 3 | 6 | 2 | 5 | 7 |
|---|---|---|---|---|---|---|

|   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|
| 0 | 1 | 1 | 0 | 0 | 1 | 1 |
| 0 | 0 | 1 | 1 | 1 | 0 | 1 |
| 1 | 1 | 0 | 1 | 0 | 1 | 1 |

# Systematic ALM Codes, $t = 1, \ell = 1$

$n = 15$  4-ary cells ( $q=4$ )

$$H = \begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 & 2 & 2 & 2 & 2 & 3 & 3 & 3 & 3 \\ 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 \end{bmatrix}$$

Encoding?

# Sufficient and Necessary Condition

$n = 15$  4-ary cells ( $q=4$ )

$$H = \begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 & 2 & 2 & 2 & 2 & 3 & 3 & 3 & 3 \\ 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 \end{bmatrix}$$

Correcting 1 error?  $\longleftrightarrow$  Unique syndromes to 1-error events

# Algebraic Condition

$n = 15$  4-ary cells ( $q=4$ )

$$H = \begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 & 2 & 2 & 2 & 2 & 3 & 3 & 3 & 3 \\ 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 \end{bmatrix}$$

Unique syndromes to  
1-error events

**NO!**  $(+1) \cdot s_1 = (+1) \cdot s_2$

$\longleftrightarrow$

**NO!**  $(+1) \cdot s_1 = \mathbf{0}$

# Systematic ALM Codes, $t = 1, \ell = 2$

$n = 12$  5-ary cells ( $q=5$ )

$$H = \begin{bmatrix} 0 & 0 & 1 & 1 & 1 & 1 & 1 & 4 & 4 & 4 & 4 & 4 \\ 1 & 4 & 0 & 1 & 2 & 3 & 4 & 0 & 1 & 2 & 3 & 4 \end{bmatrix}$$

Encoding

# Algebraic Condition, $\ell = 2$

$n = 12$  5-ary cells ( $q=5$ )

$$H = \begin{bmatrix} 0 & 0 & 1 & 1 & 1 & 1 & 1 & 4 & 4 & 4 & 4 & 4 \\ 1 & 4 & 0 & 1 & 2 & 3 & 4 & 0 & 1 & 2 & 3 & 4 \end{bmatrix}$$

**NO!**  $\alpha_1 \cdot \boxed{S_1} = \alpha_2 \cdot \boxed{S_2}$

$$\alpha_1, \alpha_2 \in \{+1, +2\}$$

# Decoding Example, $\ell = 2$

$n = 12$  5-ary cells ( $q=5$ )

$$H = \begin{bmatrix} 0 & 0 & 1 & 1 & 1 & 1 & 1 & 4 & 4 & 4 & 4 & 4 \\ 1 & 4 & 0 & 1 & 2 & 3 & 4 & 0 & 1 & 2 & 3 & 4 \end{bmatrix}$$

Received syndrome:

$$s = \begin{bmatrix} 3 \\ 1 \end{bmatrix}$$

Was it a (+1) error? No.

It was a (+2) error. Where?  $(+2)^{-1} \cdot 3 = 3 \cdot 3 = 4 \pmod{5}$

Where exactly?  $(+2)^{-1} \cdot 1 = 3 \cdot 1 = 3 \pmod{5}$

# Single Error Correction through $B_1[0, \ell]$ Sets

Definition –  $B_1[0, \ell](q)$  set:

For a set  $B$  of size  $m$ , define

$$\Omega(B) = \{ib \bmod q \mid b \in B, i \in [0, \ell]\}.$$

Then  $B$  is a  $B_1[0, \ell](q)$  set iff

$$|\Omega(B)| = \ell m + 1.$$

Examples:

$$B_1[0,1](4) = \{1,2,3\}$$

$$\Omega(B) = \{0,1,2,3\}.$$

$$B_1[0,2](5) = \{1,4\} \text{ or } \{2,3\}$$

$$\Omega(B) = \{0,1,2,3,4\}.$$



# Construction from $B_1[0, \ell]$ Sets

Examples:

1)  $B_1[0,1](4) = \{1,2,3\}$

$$H = \begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 & 2 & 2 & 2 & 2 & 3 & 3 & 3 & 3 \\ 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 \end{bmatrix}$$

2)  $B_1[0,2](5) = \{1,4\}$

$$H = \begin{bmatrix} 0 & 0 & 1 & 1 & 1 & 1 & 1 & 4 & 4 & 4 & 4 & 4 \\ 1 & 4 & 0 & 1 & 2 & 3 & 4 & 0 & 1 & 2 & 3 & 4 \end{bmatrix}$$

# 2<sup>nd</sup> Step Correctness

$$1) B_1[0,1](4) = \{1,2,3\}$$

$$H = \begin{bmatrix} 0 & 0 & 0 & 1 & 1 & 1 & 1 & 2 & 2 & 2 & 2 & 3 & 3 & 3 & 3 \\ 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 & 0 & 1 & 2 & 3 \end{bmatrix}$$

$$\alpha x_1 \neq \alpha x_2$$

# 2<sup>nd</sup> Step Correctness

$$2) B_1[0,2](5) = \{1,4\}$$

$$H = \begin{bmatrix} 0 & 0 & 1 & 1 & 1 & 1 & 1 & 4 & 4 & 4 & 4 & 4 \\ 1 & 4 & 0 & 1 & 2 & 3 & 4 & 0 & 1 & 2 & 3 & 4 \end{bmatrix}$$

$$\alpha x_1 \neq \alpha x_2$$



$$\alpha x \neq 0$$



$$\gcd(\alpha, q) = 1$$



$$\gcd(\ell!, q) = 1$$

# Perfect $B_1[0, \ell]$ Sets

$$\Omega(B) = \{ib \bmod q \mid b \in B, i \in [0, \ell]\}.$$

Upper bound:

$$|\Omega(B)| = \ell m + 1 \leq q$$

↑  
perfect if equality

$$B_1[0,1](4) = \{1,2,3\}$$

$$\Omega(B) = \{0,1,2, q - 1\}.$$

Perfect!

$$B_1[0,2](5) = \{1,4\}$$

$$\Omega(B) = \{0,1,2,3, q - 1\}.$$

Perfect!

# Extensions

- More than 1 error ( $t > 1$ )
  - $B_t[0, \ell](q)$  sets, distinct sums of t products
- Symmetric limited magnitude errors
  - $B_1[-\ell, \ell](q)$  sets
- Unbalanced limited magnitude errors
  - $B_1[-\ell_1, \ell_2](q)$  sets